

Signed Cryptographic Program Verification with Typed CRYPTOLINE

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Outline

- 1 Introduction
- 2 Previous Work & Contribution
- 3 Typed CRYPTOLINE Example
- 4 Use GCC to generate CRYPTOLINE
- 5 Case Study - NaCl
- 6 Evaluation
- 7 Conclusion

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Practical Cryptography

- Cryptographic program is written in C or ASM for **efficiency**.
- Computation over **large** finite field is not trivial in C and ASM.
- Split a large number into several smaller numbers (a.k.a. limbs).
(e.g. 4 or 5 `uint64_t`/register to store 255-bit keys for `Curve25519`)
- Computation over limbs is **error-prone**.
- A simple **bug** can cause **catastrophic** damages.
(e.g. a missing bound check in Heartbleed)



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In this work, we focus on implementation written in **C**.



Functional Correctness

So.... How to achieve the functional correctness?

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Test?

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State space is too **BIG, HARD to cover**

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Verification

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Previous Work

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- Proof Assistant + SMT Solver (CHL+14)
 - can only verify some simple code in tolerable time.
 - many human-added annotations.

SMT: Satisfiability modulo theories

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- Proof Assistant + SMT Solver + Algebra Solver (TWY17)
 - can deal with more **complex** operations like multiplication
 - SMT solver **cannot** deal with large integers multiplication well

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 - can only verify some simple code in **tolerable** time.
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 - can deal with more **complex** operations like multiplication
 - SMT solver **cannot** deal with large integers multiplication well
- **DSL + SMT Solver + Algebra Solver (PTW+18)**
 - Untyped CRYPTOLINE (only unsigned)
 - Target: ASM (some **real-word** examples in OpenSSL)
 - integer size is fixed (32/64 bit register)

SMT: Satisfiability modulo theories

DSL: Domain-specific language

Goal

- More **real-world** examples.
- Try to verify the C implementation **once** instead of ASM for every platforms.
 - most implementation now are still written in C instead of human-optimized ASM
- Less verification effort and friendly to normal cryptographic library developers.

Target Cryptographic Libraries

- OpenSSL: **UBIQUITOUS**
- BoringSSL: Chrome, Android
- NaCl: reference implementation
- wolfSSL: embedded systems
- Bitcoin's libsecp256k1: ECDSA used by **MANY** cryptocurrencies
(Ethereum, Zcash, Ripple, ...)

What Curves We Verified

- OpenSSL:
 - NIST P-224 : $2^{224} - 2^{96} + 1$ 32/64: integer size (unsigned 64)
 - NIST P-256 : $2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$ (unsigned 64)
 - NIST P-521 : $2^{521} - 1$ (unsigned 64)
 - Curve25519 : $2^{255} - 19$ (unsigned 64, **signed** 32)
- BoringSSL: Curve25519 (unsigned 64)
- NaCl: Curve25519 (unsigned 64, **signed** 64)
- wolfSSL: Curve25519 (same as OpenSSL's) (**signed** 32)
- Bitcoin: Secp256k1 ($2^{256} - 2^{32} - 2^9 - 2^8 - 2^7 - 2^6 - 2^4 - 1$) (unsigned, **signed**)

Contribution

- **Typed CRYPTOLINE** – unsigned and **signed, arbitrary size integers**
 - type system (type checking & type inference)
- A **GCC** plugin that translates **GIMPLECRYPTOLINE** into **Typed CRYPTOLINE**
- **GIMPLECRYPTOLINE** – a subset of **GIMPLE**
 - **GIMPLE**: a GCC IR used in **machine-independent optimization**
- Verify **GIMPLE** code after **machine-independent optimization**
- First to verify **signed C** implementation in cryptographic libraries used in **industry**
- Found a **bug** in NaCl's Curve25519 - Case study

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Typed CRYPTOLINE Program

- Program - instructions
- Specification
 - Assumption (Precondition)
 - Assertion (Postcondition)
 - Properties {algebra && range}
 - range: variables should be in a **proper** range (e.g. $a < 2^{51}$)
checked by SMT solver (Boolector, MathSAT, Z3 ...)
 - algebra: mathematical properties (e.g. $c = a \times b$)
checked by algebraic solver (Sage, Singular, Mathematica ...)
- Hoare triple: {assumption} program {assertion}

Typed CRYPTOLINE Program Example - Naive Addition

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1 proc main (uint64 a0, uint64 a1, uint64 b0, uint64 b1) =
2 {
3     true // algebraic prop; true means no assumption
4     &&
5     and [ // range prop
6         a0 <u (2**63)@64, a1 <u (2**63)@64,
7         b0 <u (2**63)@64, b1 <u (2**63)@64
8     ]
9 }
10 add c0 a0 b0; // c0 = a0 + b0
11 add c1 a1 b1; // c1 = a1 + b1
12 {
13     limbs 64 [c0, c1]
14 =
15     limbs 64 [a0, a1] + limbs 64 [b0, b1]
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$$\text{limbs } 64 [a_0, a_1, \dots, a_n] = \sum_{i=0}^n a_i \times 2^{64 \times i}$$

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$$2^{63} - 1 + 2^{63} - 1 = 2^{64} - 2 \leq 2^{64} - 1$$

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$$2^{63} + 2^{63} = 2^{64} \not\leq 2^{64} - 1$$
$$2^{64} = 0 \pmod{2^{64}}$$

Program Safety Check by SMT Solver

Safety in our context means that following kinds of errors do not exist.

- Overflow / Underflow

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- Overflow / Underflow
- Cast between types
(`uint64 ↔ int64`, `uint64 ↔ uint32`)
- Value preserving casting (vpc)

2's complement representation for signed integers

`uint4 ↔ int4`

$(0111)_2 = 7$ (unsigned) = 7 (signed)
 $(1111)_2 = 15$ (unsigned) = -1 (signed)

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		<code>uint4 ↔ int4</code>		
$(0111)_2 =$	7	(unsigned) =	7	(signed) ✓
$(1111)_2 =$	15	(unsigned) =	-1	(signed) ✗

BUG (vpc) or **on purpose** (cast)

Counterexample by SMT solver

Typed CRYPTOLINE Program Example - Cast v.s. VPC

```
1 proc main (uint64 a ,uint64 b)= 1 proc main (uint64 a ,uint64 b)=  
2 { 2 {  
3   true 3   true  
4   && 4   &&  
5   and [ 5   and [  
6     a <u (2**63), b <u (2**63) 6     a <u (2**63), b <u (2**63)  
7   ] 7   ]  
8 } 8 }  
9 cast wa@int64 a; 9 vpc wa@int64 a;   
10 cast wb@int64 b; 10 vpc wb@int64 b;   
11 mul c wa wb; 11 mul c wa wb;  
12 { ... } 12 { ... }
```

Figure: cast = vpc in some cases

under the assumption, sign bit will **never** be 1.

Typed CRYPTOLINE Program Example - VPC Error

```
1 proc main (uint64 a ,uint64 b)= 1 proc main (uint64 a ,uint64 b)=  
2 { 2 {  
3 true 3 true  
4 && 4 &&  
5 and [ 5 and [  
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Figure: cast \neq vpc in some cases

$$2^{63} = (100\ldots0)_2$$

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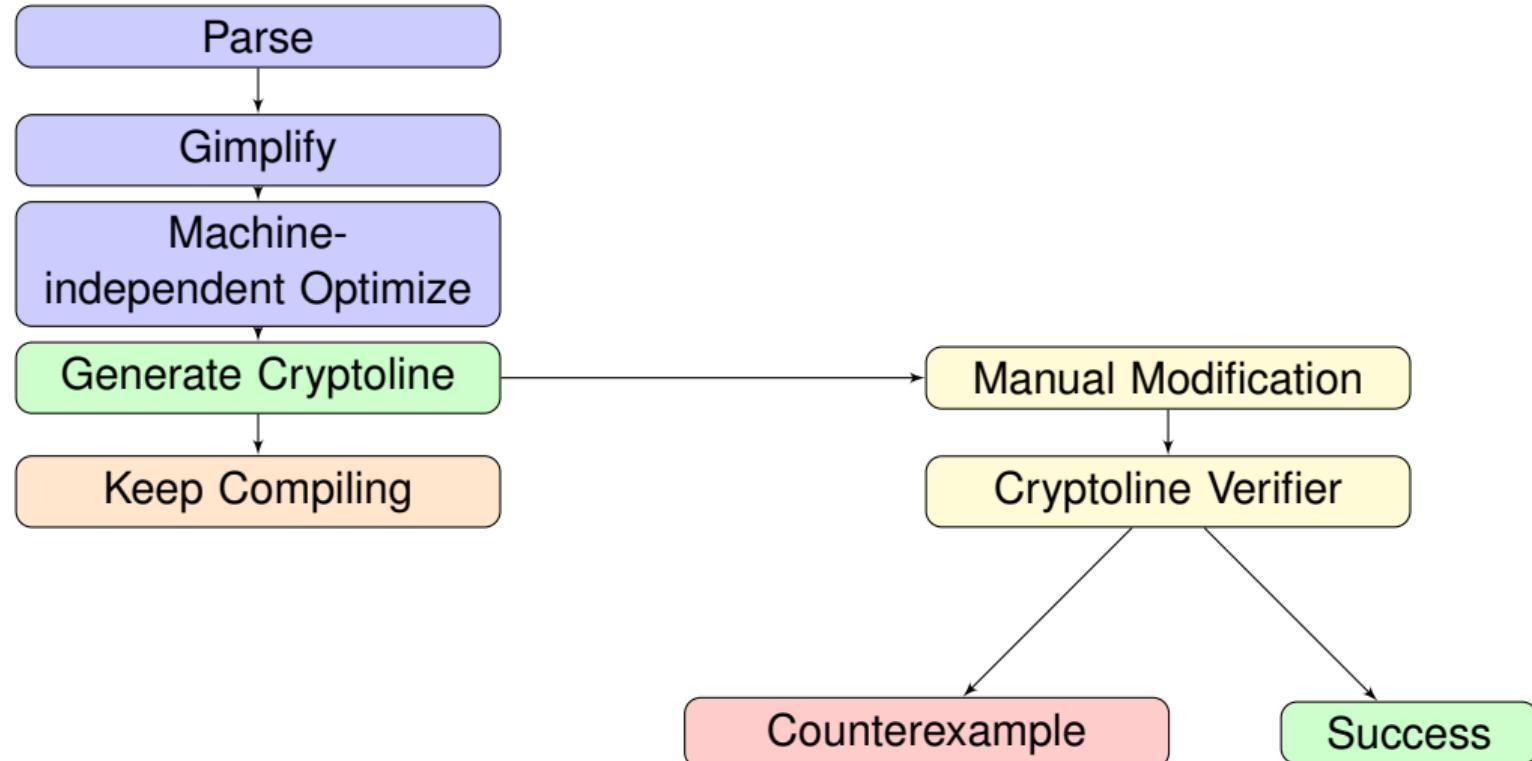
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GCC Plugin

- Introduced in GCC 4.5.0
- Let us add **custom** optimization passes
- Able to access **AST** (abstract syntax tree)
 - No need to write parser by yourself!



Verification Workflow Using GCC Plugin



GIMPLE Example

```
1  f0_3 = *f_2(D);
2  f1_4 = MEM[ (const int32_t*) f_2(D) +4B];
3  ...
4  g0_14 = *g_13(D);
5  g1_15 = MEM[ (const int32_t*) g_13(D) +4B];
6  ...
7  h0_24 = f0_3 - g0_14;
8  h1_25 = f1_4 - g1_15;
9  ...
10 *h_34(D) = h0_24;
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12 ...
```

?LHS = MEM[?RHS] \Rightarrow Load from RHS to LHS
MEM[?LHS] = ?RHS \Rightarrow Store RHS to LHS

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GIMPLECRYPTOLINE Example

generated by the plugin **automatically**. later manually add assumption / assertion.

```
1 proc main () =
2 { true && true }
3 mov f03 f2_0; // f0_3 = *f_2
4 mov f14 f2_4; // f1_4 = MEM[ (...) f_2 + 4]
5 ...
6 mov g014 g13_0;
7 mov g115 g13_4;
8 ...
9 sub h024 f03 g014; // h0_24 = f0_3 - g0_14
10 sub h125 f14 g115;
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Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;
/* Find the difference of two numbers: output = in - output
 * (note the order of the arguments!)
 */
static void fdifference_backwards(felem *ioutput, const felem *iin) {
    static const int64_t twotothe51 = (1L << 51);
    const int64_t *in = (const int64_t *) iin;
    int64_t *out = (int64_t *) ioutput;

    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];
    out[4] = in[4] - out[4];

    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
    NEGCHAIN19(4, 0);
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;  
/* Find the difference of two numbers: output = in - output  
 * (note the order of the arguments!)  
 */  
static void fdifference_backwards(felem *ioutput, const felem *iin) {  
    static const int64_t twotothe51 = (1L << 51);  
    const int64_t *in = (const int64_t *) iin;  
    int64_t *out = (int64_t *) ioutput;  
  
    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];  
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];  
    out[4] = in[4] - out[4];  
  
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);  
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);  
    NEGCHAIN19(4, 0);  
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);  
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;
/* Find the difference of two numbers: output = in - output
 * (note the order of the arguments!)
 */
static void fdifference_backwards(felem *ioutput, const felem *iin) {
    static const int64_t twotothe51 = (1L << 51);
    const int64_t *in = (const int64_t *) iin;
    int64_t *out = (int64_t *) ioutput;

    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];
    out[4] = in[4] - out[4];

    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
    NEGCHAIN19(4, 0);
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;
/* Find the difference of two numbers: output = in - output
 * (note the order of the arguments!)
 */
static void fdifference_backwards(felem *ioutput, const felem *iin) {
    static const int64_t twotothe51 = (1L << 51);
    const int64_t *in = (const int64_t *) iin;
    int64_t *out = (int64_t *) ioutput;

    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];
    out[4] = in[4] - out[4];

    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
    NEGCHAIN19(4, 0);
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;
/* Find the difference of two numbers: output = in - output
 * (note the order of the arguments!)
 */
static void fdifference_backwards(felem *ioutput, const felem *iin) {
    static const int64_t twotothe51 = (1L << 51);
    const int64_t *in = (const int64_t *) iin;
    int64_t *out = (int64_t *) ioutput;

    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];
    out[4] = in[4] - out[4];

    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
    NEGCHAIN19(4, 0);
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

Case Study - Field Subtraction in NaCl Curve25519

```
typedef uint64_t felem;
/* Find the difference of two numbers: output = in - output
 * (note the order of the arguments!)
 */
static void fdifference_backwards(felem *ioutput, const felem *iin) {
    static const int64_t twotothe51 = (1L << 51);
    const int64_t *in = (const int64_t *) iin;
    int64_t *out = (int64_t *) ioutput;

    out[0] = in[0] - out[0]; out[1] = in[1] - out[1];
    out[2] = in[2] - out[2]; out[3] = in[3] - out[3];
    out[4] = in[4] - out[4];

    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
    NEGCHAIN19(4, 0);
    NEGCHAIN(0, 1); NEGCHAIN(1, 2);
    NEGCHAIN(2, 3); NEGCHAIN(3, 4);
```

5 `uint64` limbs and use `signed` computation

```
| int64_t t;  
|  
#define NEGCHAIN(a,b) \  
t = out[a] >> 63; \  
out[a] += twotothe51 & t; \  
out[b] -= 1 & t;  
  
#define NEGCHAIN19(a,b) \  
t = out[a] >> 63; \  
out[a] += twotothe51 & t; \  
out[b] -= 19 & t;  
  
NEGCHAIN(0, 1);  
NEGCHAIN(1, 2);  
NEGCHAIN(2, 3);  
NEGCHAIN(3, 4);  
NEGCHAIN19(4, 0);  
NEGCHAIN(0, 1);  
NEGCHAIN(1, 2);  
NEGCHAIN(2, 3);  
NEGCHAIN(3, 4);  
}
```

Figure: Bitwise tricks (signed right shift) & Reduction chain

```
| int64_t t;
```

```
#define NEGCHAIN(a,b) \
    t = out[a] >> 63; \
    out[a] += twotothe51 & t; \
    out[b] -= 1 & t;
```

```
#define NEGCHAIN19(a,b) \
    t = out[a] >> 63; \
    out[a] += twotothe51 & t; \
    out[b] -= 19 & t;
```

```
NEGCHAIN(0, 1);  
NEGCHAIN(1, 2);  
NEGCHAIN(2, 3);  
NEGCHAIN(3, 4);  
NEGCHAIN19(4, 0);  
NEGCHAIN(0, 1);  
NEGCHAIN(1, 2);  
NEGCHAIN(2, 3);  
NEGCHAIN(3, 4);  
}
```

Figure: Bitwise tricks (signed right shift) & Reduction chain

sign_bit(out[a]) == 1/0 \leftrightarrow t is all 1/0

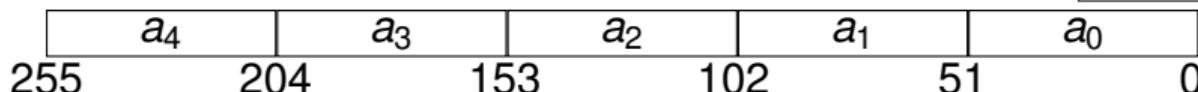
Verify by CRYPTOLINE - Pre-condition

```
proc main (uint64 a0, uint64 a1, uint64 a2, uint64 a3, uint64 a4,  
          uint64 b0, uint64 b1, uint64 b2, uint64 b3, uint64 b4) =  
{  
    true  
    &&  
    and [  
        a0 <u (2**51)@64,  
        a1 <u (2**51)@64,  
        a2 <u (2**51)@64,  
        a3 <u (2**51)@64,  
        a4 <u (2**51)@64,  
        b0 <u (2**51)@64,  
        b1 <u (2**51)@64,  
        b2 <u (2**51)@64,  
        b3 <u (2**51)@64,  
        b4 <u (2**51)@64  
    ]  
}
```

Assume

Program

Assert



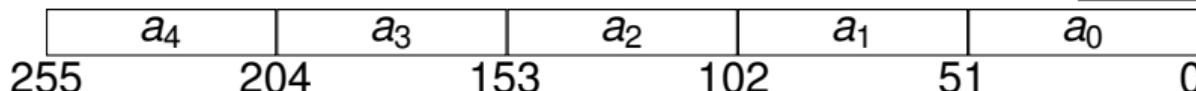
Verify by CRYPTOLINE - Pre-condition

```
proc main (uint64 a0, uint64 a1, uint64 a2, uint64 a3, uint64 a4,  
          uint64 b0, uint64 b1, uint64 b2, uint64 b3, uint64 b4) =  
{  
    true  
    &&  
    and [  
        a0 <u (2**51)@64,  
        a1 <u (2**51)@64,  
        a2 <u (2**51)@64,  
        a3 <u (2**51)@64,  
        a4 <u (2**51)@64,  
        b0 <u (2**51)@64,  
        b1 <u (2**51)@64,  
        b2 <u (2**51)@64,  
        b3 <u (2**51)@64,  
        b4 <u (2**51)@64  
    ]  
}
```

Assume

Program

Assert



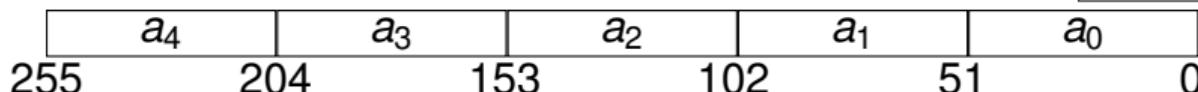
Verify by CRYPTOLINE - Pre-condition

```
proc main (uint64 a0, uint64 a1, uint64 a2, uint64 a3, uint64 a4,
           uint64 b0, uint64 b1, uint64 b2, uint64 b3, uint64 b4) =
{
    true
    &&
    and [
        a0 <u (2**51)@64,
        a1 <u (2**51)@64,
        a2 <u (2**51)@64,
        a3 <u (2**51)@64,
        a4 <u (2**51)@64,
        b0 <u (2**51)@64,
        b1 <u (2**51)@64,
        b2 <u (2**51)@64,
        b3 <u (2**51)@64,
        b4 <u (2**51)@64
    ]
}
```

Assume

Program

Assert



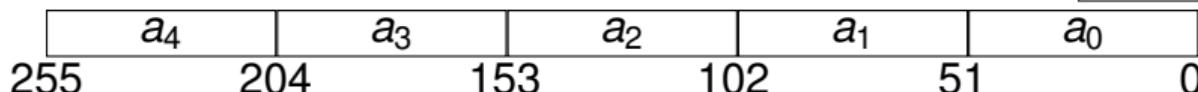
Verify by CRYPTOLINE - Pre-condition

```
proc main (uint64 a0, uint64 a1, uint64 a2, uint64 a3, uint64 a4,
           uint64 b0, uint64 b1, uint64 b2, uint64 b3, uint64 b4) =
{
    true
    &&
    and [
        a0 <u (2**51)@64,
        a1 <u (2**51)@64,
        a2 <u (2**51)@64,
        a3 <u (2**51)@64,
        a4 <u (2**51)@64,
        b0 <u (2**51)@64,
        b1 <u (2**51)@64,
        b2 <u (2**51)@64,
        b3 <u (2**51)@64,
        b4 <u (2**51)@64
    ]
}
```

Assume

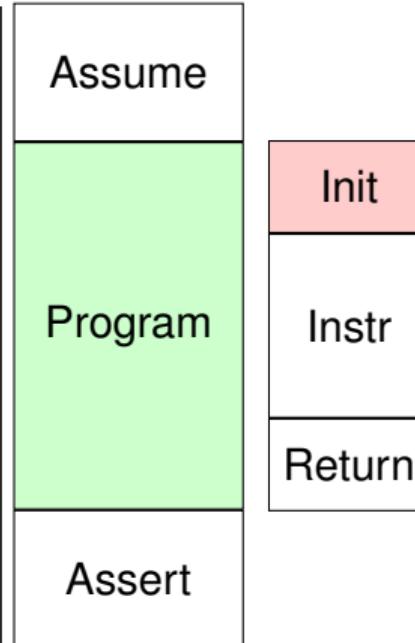
Program

Assert



Verify by CRYPTOLINE - Init type casting

```
vpc iin52_0@int64 a0;  
vpc iin52_8@int64 a1;  
vpc iin52_16@int64 a2;  
vpc iin52_24@int64 a3;  
vpc iin52_32@int64 a4;  
vpc ioutput53_0@int64 b0;  
vpc ioutput53_8@int64 b1;  
vpc ioutput53_16@int64 b2;  
vpc ioutput53_24@int64 b3;  
vpc ioutput53_32@int64 b4;
```



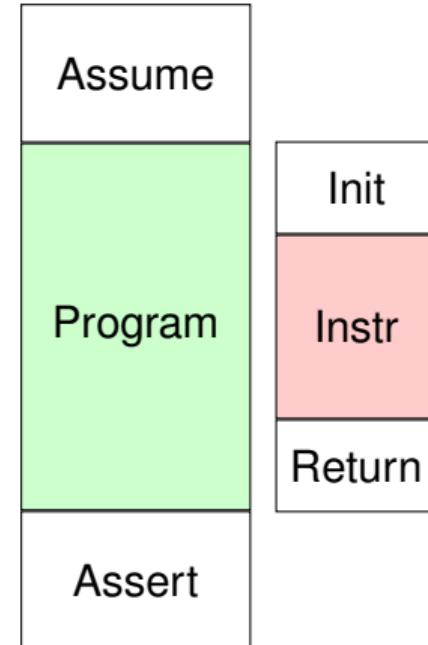
$\text{uint64} \rightarrow \text{int64}$

bridge **assumption** and program

vpc: value preserve casting (will do safety check)

Verify by CRYPTOLINE - Instructions

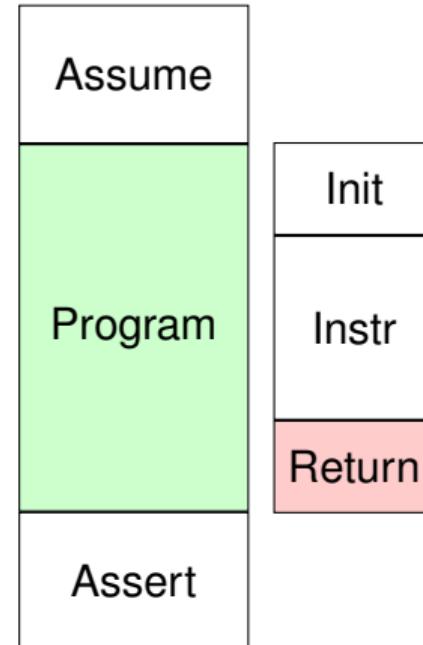
```
(* _1 = MEM[(const int64_t * )iin_52(D)]; *)
mov v1 iin52_0;
(* _2 = MEM[(int64_t * )ioutput_53(D)]; *)
mov v2 ioutput53_0;
(* _3 = _1 - _2; *)
ssub v3 v1 v2;
(* MEM[(int64_t * )ioutput_53(D)] = _3; *)
mov ioutput53_0 v3;
(* _4 = MEM[(const int64_t * )iin_52(D) + 8B]; *)
mov v4 iin52_8;
(* _5 = MEM[(int64_t * )ioutput_53(D) + 8B]; *)
mov v5 ioutput53_8;
(* _6 = _4 - _5; *)
ssub v6 v4 v5;
```



ssub: signed subtraction (usub/ssub explicitly \Rightarrow type checking, sub \Rightarrow type inference)

Verify by CRYPTOLINE - Return type casting

```
vpc c0@uint64 ioutput53_0@int64;
vpc c1@uint64 ioutput53_8@int64;
vpc c2@uint64 ioutput53_16@int64;
vpc c3@uint64 ioutput53_24@int64;
vpc c4@uint64 ioutput53_32@int64;
```



$\text{int64} \rightarrow \text{uint64}$

bridge program and assertion

vpc: value preserve casting (will do safety check)

Verify by CRYPTOLINE - Post-condition

```
{  
    limbs 51 [c0, c1, c2, c3, c4])  
    =  
    (  
        limbs 51 [a0, a1, a2, a3, a4])  
        -  
        (limbs 51 [b0, b1, b2, b3, b4])  
    )  
    (mod (2**255 - 19))  
    &&  
    and [  
        c0 <u (2**51)@64,  
        c1 <u (2**51)@64,  
        c2 <u (2**51)@64,  
        c3 <u (2**51)@64,  
        c4 <u (2**51)@64  
    ]  
}
```

Assume

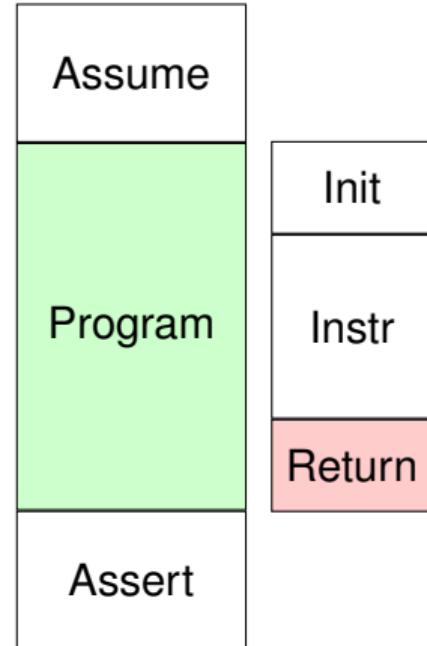
Program

Assert

$$\text{limbs } 51[a_0, a_1, \dots, a_n] = \sum_{i=0}^n a_i \times 2^{51 \times i} \\ \text{mod m: under modulo m}$$

Verify by CRYPTOLINE - Return type casting - Revisit

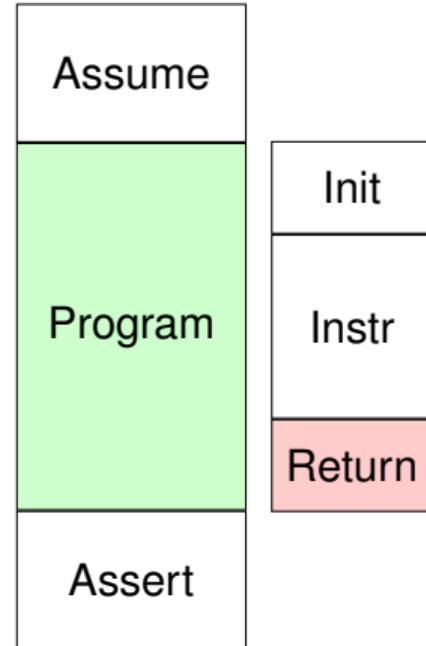
```
vpc c0@uint64 ioutput53_0@int64;  
vpc c1@uint64 ioutput53_8@int64;  
vpc c2@uint64 ioutput53_16@int64;  
vpc c3@uint64 ioutput53_24@int64;  
vpc c4@uint64 ioutput53_32@int64;
```



vpc: value preserve casting (will do safety check)

Verify by CRYPTOLINE - Return type casting - Revisit

```
vpc c0@uint64 ioutput53_0@int64;  
vpc c1@uint64 ioutput53_8@int64;  
vpc c2@uint64 ioutput53_16@int64;  
vpc c3@uint64 ioutput53_24@int64;  
vpc c4@uint64 ioutput53_32@int64;
```



vpc: value preserve casting (will do safety check)

Safety Check Failed

Counterexample generated by SMT solvers

```
(b4_0 (_ bv0 64))  
(b3_0 (_ bv0 64))  
(b2_0 (_ bv2251799813685250 64))  
(b1_0 (_ bv0 64))  
(b0_0 (_ bv2 64))  
(a4_0 (_ bv0 64))  
(a3_0 (_ bv1 64))  
(a2_0 (_ bv0 64))  
(a1_0 (_ bv1 64))  
(a0_0 (_ bv2 64)) )
```

Figure: output by MathSAT

Counterexample generated by SMT solvers

```
(b4_0 (_ bv0 64))
(b3_0 (_ bv0 64))
(b2_0 (_ bv2251799813685250 64))
(b1_0 (_ bv0 64))
(b0_0 (_ bv2 64))
(a4_0 (_ bv0 64))
(a3_0 (_ bv1 64))
(a2_0 (_ bv0 64))
(a1_0 (_ bv1 64))
(a0_0 (_ bv2 64)) )
```

Figure: output by MathSAT

Found Counterexample translated to C language

```
int main()
{
    felem in[5] = { 2, 1, 0, 1, 0 };
    felem out[5] = { 2, 0, 2251799813685250, 0, 0 };
    fdifference_backwards(out, in);
    for (int i = 0; i < 5; i++) {
        printf(" out%d: 0x%llx \n", i, out[i]);
    }
}
```

Check whether the program result is correct !

out0: 0x0

out1: 0x1

out2: 0x7fffffff

out3: 0x7fffffff

out4: 0xffffffff

Check whether the program result is correct !

out0:	0x0	0x0
out1:	0x1	0x1
out2:	0x7fffffffffffffe	0x5f6080e
out3:	0x7fffffff00000000	0x0
out4:	0xffffffffffffffffffff	0x0

Underflow and not in proper range

$$0xffffffffffffffffffff = 2^{64} - 1$$

Outline

- 1 Introduction
- 2 Previous Work & Contribution
- 3 Typed CRYPTOLINE Example
- 4 Use GCC to generate CRYPTOLINE
- 5 Case Study - NaCl
- 6 Evaluation
- 7 Conclusion

Glimpse Evaluation Result

- 82 C functions (when paper is submitted)
- Evaluated on two different machines
 - much more range properties and safety check by SMT solver ⇒ done in parallel
 - a few algebraic properties (most have only 1)
 - field operation
 - group operation

M1: Macbook 13" 2C/4T 16GB

M2: Ubuntu Server 18C/36T 1024GB

Evaluation Table - all functions

Table 2: Experimental Result

Function	L_{IR}	L_{CL}	D	P	TR_{M1}	MR_{M1}	TA_{M1}	MAM_1	TR_{M2}	MR_{M2}	TA_{M2}	MAM_2
nac1/curve25519/donna_c64/curve25519.c (MathSAT, SMT-LIB2 format)												
difference_backwards	69	69	66	0	-	-	0.23	6.3	-	-	0.14	9.1
fmul	91	127	10	14	12.51	452.2	0.20	6.3	4.05	486.6	0.14	9.5
fscalar_product	58	38	7	10	27.5	104.4	0.20	5.6	0.95	108.4	0.12	8.6
fsquare	68	116	10	12	7.44	288.1	0.22	6.3	2.61	301.0	0.13	9.3
fsum	20	20	0	0	0.48	5.6	0.15	4.8	0.22	10.0	0.10	8.2
fmonty	1147	1493	361	127	-	-	OOM	OOM	-	-	353.66	32764
wolfssl/fe_operations.c (Booletoor with Lingeling, BTOR format)												
fe_add	40	40	0	0	1.48	6.5	0.19	5.6	0.61	9.5	0.11	8.6
fe_mul	305	305	20	24	OOT	OOT	0.32	7.0	13178	883.3	0.15	9.9
fe_mull21666	91	91	20	20	19.68	17.9	0.26	6.4	3.75	13.8	0.13	9.4
fe_neg	30	30	0	0	1.24	6.5	0.18	5.3	0.63	9.3	0.10	8.3
fe_sq	204	204	20	24	13411.84	351.9	0.33	6.7	2033	355.6	0.14	9.6
fe_sq2	214	214	20	24	18252.02	388.9	0.30	6.8	2763	385.5	0.14	9.6
fe_sub	40	40	0	0	1.31	6.5	0.16	5.7	0.64	9.4	0.11	8.6
curve25519	2770	2770	200	236	OOT	OOT	12.06	385.6	68140	796.7	8.26	382.1
bitcoinc/field_5x52_impl.h (MathSAT, SMT-LIB2 format)												
secp256k1_fe_add	13	20	0	0	0.33	5.3	0.14	4.8	0.22	10.0	0.09	8.3
secp256k1_fe_cmov	29	49	13	20	1.35	28.7	0.29	6.4	0.46	29.6	0.17	9.3
secp256k1_fe_from_storage	24	32	6	14	0.53	6.4	0.15	5.2	0.31	10.7	0.09	8.4
secp256k1_fe_mull_int	16	16	2	0	0.52	26.1	0.14	4.7	0.28	28.0	0.10	8.4
secp256k1_fe_negate	20	20	2	0	0.52	5.7	0.18	4.9	0.27	9.9	0.11	8.6
bitcoinc/field_5x52_impl.h (Booletoor with Lingeling, BTOR format)												
secp256k1_fe_normalize	52	60	21	0	117.18	45.3	0.12	5.3	91.89	31.5	0.08	8.3
secp256k1_fe_normalize_var	63	63	29	0	120.80	47.1	0.12	5.4	95.65	34.1	0.08	8.3
secp256k1_fe_normalize_weak	26	26	15	0	63.83	40.0	0.25	5.3	51.51	28.3	0.13	8.8
secp256k1_fe_normalize_to_zero	34	39	10	0	203.12	60.3	0.16	5.2	151.03	42.9	0.08	8.2
bitcoinc/field_5x52_impl.h (MathSAT, SMT-LIB2 format)												
secp256k1_fe_mul_inner	111	137	17	24	16.09	46.0	0.22	6.5	4.00	489.1	0.14	9.5
secp256k1_fe_sqr_inner	90	116	21	22	9.91	284.5	0.20	6.4	2.72	303.2	0.14	9.3
bitcoinc/scalar_4x64_impl.h (MathSAT, SMT-LIB2 format)												
secp256k1_scalar_addr	81	102	55	22	2.03	10.1	0.21	6.5	1.11	14.1	0.13	9.4
secp256k1_scalar_eq	17	17	23	0	0.29	9.2	0.10	4.7	0.26	14.5	0.07	7.6
secp256k1_scalar_mul_512	273	384	136	90	13.73	263.3	0.26	7.1	4.9	280.0	0.16	9.9
secp256k1_scalar_mul	652	947	379	228	128.19	453.9	0.84	19.8	741.35	2219	0.43	16.3
secp256k1_scalar_reduce	41	55	4	1	28.50	132.4	0.10	5.6	40.31	135.5	0.08	8.0
secp256k1_scalar_reduce_512	379	563	243	138	31.84	127.5	0.37	8.7	8.25	128.2	0.23	11.7
secp256k1_scalar_reduce	34	32	11	8	1.52	11.7	0.18	6.4	0.88	15.2	0.14	9.3
secp256k1_scalar_sqrt_512	235	333	145	88	23.75	212.9	0.26	7.2	7.39	204.8	0.17	10.1
secp256k1_scalar_sqrt	614	896	388	226	234.87	349.1	0.82	19.8	26.69	341.5	0.45	16.5
bitcoinc/group_impl.h (MathSAT, SMT-LIB2 format)												
secp256k1_ge_from_storage	48	65	12	28	0.93	6.5	0.19	6.3	0.48	10.7	0.12	9.2
secp256k1_ge_neg	33	31	0	10	0.76	6.6	0.19	5.4	0.44	11.2	0.13	8.7
secp256k1_ge_add_ge_var	2109	2437	371	396	574.39	3166.9	OOT	OOT	75	3344	9363	70156
secp256k1_ge_double_var	899	1042	154	160	163.30	170.93	0.77	18.4	25.27	1800	0.57	22.7
openssl/curve25519.e (MathSAT, SMT-LIB2 format)												
fe51_add	20	20	0	0	0.85	6.0	0.19	4.9	0.36	10.0	0.10	8.3
fe51_mul	96	105	11	20	17.95	385.2	0.26	6.4	3.69	409.3	0.13	9.2
fe51_mull21666	44	44	11	14	1.3	17.3	0.25	5.6	0.63	20.2	0.12	8.7
fe51_sq	73	82	11	0	8.07	227.0	0.23	6.3	2.22	247.6	0.14	9.2

Function	L_{IR}	L_{CL}	D	P	TR_{M1}	MR_{M1}	TA_{M1}	MA_{M1}	TR_{M2}	MR_{M2}	TA_{M2}	MA_{M2}
fe51_sub	25	25	10	10	0.37	6.8	0.24	5.4	0.26	11.4	0.13	8.9
x25519_scalar_mult	923	1047	110	194	558.56	1419.8	187.40	5538	119.89	1472	145.12	5511
openssl/ecc_nistp224.c (MathSAT, SMT-LIB2 format)												
felem_diff_128_64	24	36	0	0	0.56	6.4	0.23	5.1	0.32	10.7	0.14	8.6
felem_diff	24	24	0	0	0.55	5.8	0.19	4.9	0.33	10.4	0.11	8.8
felem_mul	40	40	0	0	2.24	83.2	0.15	5.2	0.65	88	0.09	8.2
felem_mull_reduce	82	121	15	16	10.65	321.8	0.20	6.4	3.11	322.5	0.13	9.1
felem_neg	47	58	5	10	0.95	6.8	0.19	5.8	0.55	11.1	0.12	8.7
felem_reduce	56	95	6	18	1.67	13.7	0.20	6.3	0.88	17.3	0.13	9.3
felem_scalar	12	12	0	0	0.48	26.7	0.14	4.6	0.24	28.9	0.09	8.1
felem_square	27	27	0	0	1.11	45.1	0.15	4.9	0.43	47.6	0.10	8.2
felem_square_reduce	69	108	14	18	6.36	195.8	0.21	6.4	1.81	198.8	0.13	9.2
felem_sum	16	16	0	0	0.41	5.4	0.15	4.7	0.26	10.0	0.10	8.3
widefelem_diff	41	63	0	0	0.90	6.5	0.19	5.7	0.46	10.6	0.12	8.7
widfefelem_scalar	21	21	0	0	2.58	87.7	0.14	4.8	0.70	88.3	0.10	8.4
openssl/ecc_nistp256c (MathSAT, SMT-LIB2 format)												
felem_diff	24	36	0	0	0.59	7.6	0.18	5.1	0.35	11.7	0.12	8.6
felem_scalar	13	13	0	0	0.70	47.7	0.17	4.6	0.31	48.8	0.10	8.2
felem_shrink	65	95	18	16	1.78	14.0	0.20	6.4	0.95	17.1	0.13	9.3
felem_small_mmul	145	95	17	46	4.75	123.0	0.23	7.0	2.29	123.2	0.14	9.8
felem_small_sum	20	20	0	0	0.41	5.8	0.14	4.8	0.25	10.2	0.10	8.4
felem_sum	16	16	0	0	0.41	5.6	0.14	4.7	0.24	10.3	0.09	8.2
smallfelem_mul	88	136	0	30	2.80	91.9	0.17	6.4	1.22	95.4	0.11	9.4
smallfelem_neg	26	26	0	0	0.1	5.4	0.19	4.9	0.27	9.7	0.12	8.6
smallfelem_square	60	108	0	20	1.92	55.8	0.15	6.3	0.85	55.5	0.10	9.2
openssl/ecc_nistp521c (MathSAT, SMT-LIB2 format)												
felem_diff64	45	45	18	18	0.81	6.9	0.20	6.4	0.48	11.4	0.13	9.3
felem_diff128	45	72	18	18	1.13	7.9	0.21	6.4	0.47	11.9	0.12	9.2
felem_neg	27	27	0	0	0.77	6.4	0.18	5.3	0.48	10.0	0.12	8.6
felem_reduce	122	155	74	72	4.10	7.8	0.24	6.7	2.06	10.8	0.14	9.6
felem_scalar	27	27	0	0	0.80	28.4	0.14	5.0	0.36	29.0	0.09	8.3
felem_scalar64	27	27	0	0	0.82	28.2	0.15	4.9	0.35	28.9	0.09	8.3
felem_scalar128	27	27	0	0	1.26	48.4	0.14	5.0	0.41	48.8	0.09	8.4
felem_sum64	36	36	0	0	0.49	6.0	0.14	5.2	0.29	10.0	0.10	8.3
felem_diff_128_64	54	54	0	0	1.34	7.2	0.29	6.0	0.68	11.4	0.15	8.7
felem_mul	188	188	0	0	23.92	187.0	0.22	6.6	3.13	182.5	0.13	9.5
felem_square	111	111	0	0	7.38	95.5	0.21	6.4	0.99	103.9	0.13	9.3
boringssl/fiat/curve25519.c (MathSAT, SMT-LIB2 format)												
fe_add	11	20	0	0	0.33	5.3	0.14	4.8	0.20	10.0	0.10	8.2
fe_mul_impl	96	108	9	22	18.39	452.9	0.21	6.4	5.11	473.9	0.13	9.2
fe_mul21666	43	43	9	14	1.12	18.4	0.20	5.7	0.62	21.2	0.11	8.6
fe_sqr_Impl	73	85	9	22	10.59	278.7	0.26	6.3	3.11	293.0	0.12	9.2
fe_sub	15	25	0	0	0.51	5.9	0.19	5.0	0.28	10.4	0.11	8.8
x25519_scalar_mult_generic	927	1073	161	212	470.68	1489.0	120.33	5726	118.95	1579	91.99	5766

Some comparisons

Montgomery Ladder step* involves 4 add, 4 sub, 4 square, 6 mul
(Curve25519) (field operations)

F	U/S	L_{IR}	L_{CL}	TR_{M1}	TA_{M1}	TR_{M2}	TA_{M2}	TH
openSSL	5 * U64	923	1047	9.3m	0.93s	2m	0.61s	
boringSSL	5 * U64	927	1073	7.8m	0.89s	2m	0.56s	
boringSSL	10 * U32	2715	3419	27.5m	59s	6.3m	42s	2h
wolfSSL	10 * S32	2770	2770	OOT	12s	18.9h	8s	

L_{IR} : lines of IR

L_{CL} : lines of CRYPTOLINE

TR(range, safety), TA(algebra): used time

OOT: used time > 1day

TH: human effort (one person)

Montgomery Ladder is used for scalar multiplication of elliptic curve point

$$Q = aP$$

Outline

- 1 Introduction
- 2 Previous Work & Contribution
- 3 Typed CRYPTOLINE Example
- 4 Use GCC to generate CRYPTOLINE
- 5 Case Study - NaCl
- 6 Evaluation
- 7 Conclusion

Conclusion

- A lightweight and easy to use method to verify cryptographic software involving both unsigned/signed operations.
- A GCC Plugin reducing human effort
- Verify several functions in well-known cryptographic libraries.
 - OpenSSL
 - BoringSSL
 - NaCl
 - wolfSSL
 - Bitcoin's libsecp256k1



CryptoLine Verifier



GCC Plugin



This Slide¹



github.com/fmlab-iis

Signed Cryptographic Program Verification with Typed CRYPTO_{LINE}
[Open Access](#)

¹twleo.com/slides/ccs19-slide.pdf